

NanoCF:

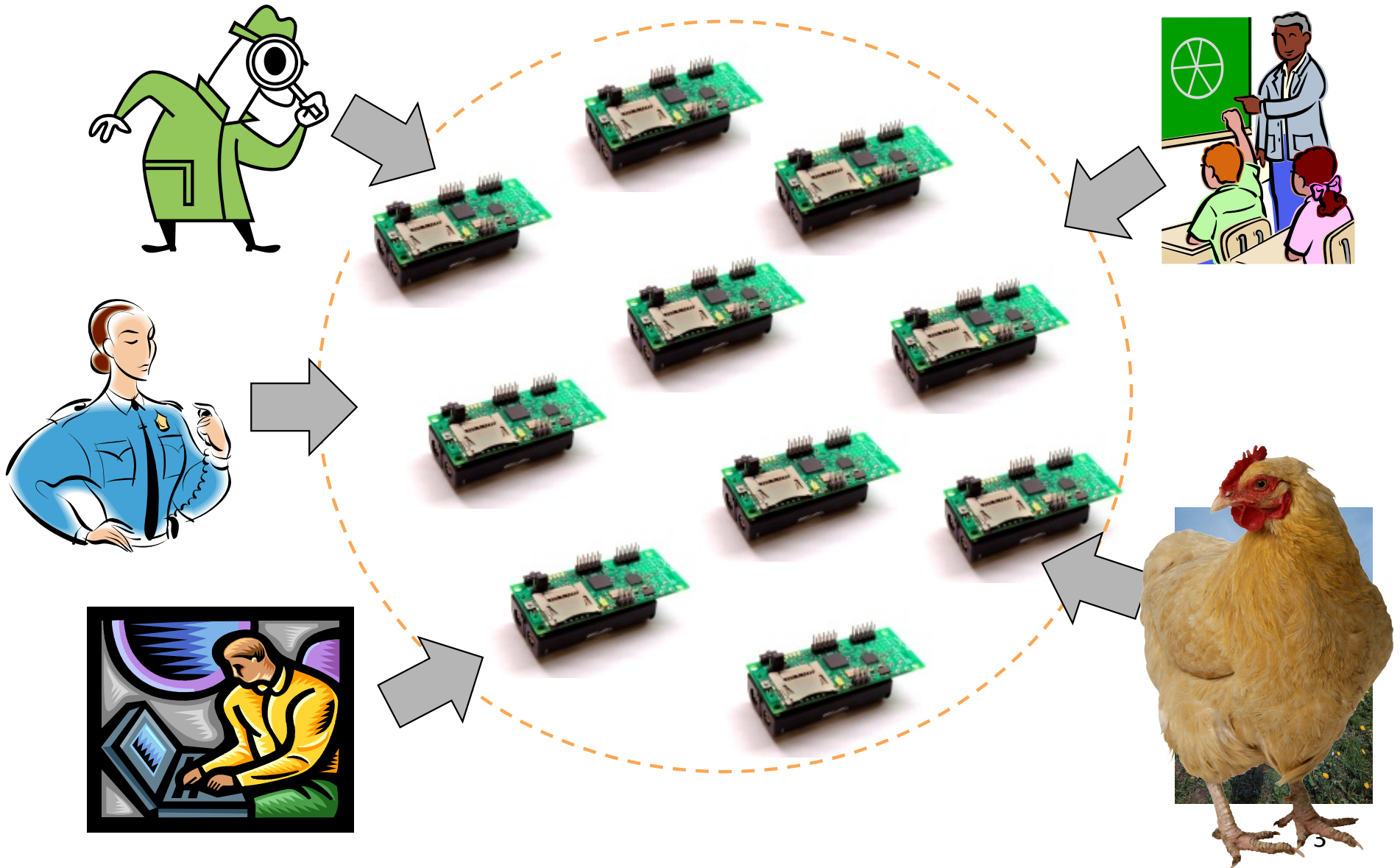
A Coordination Framework for Multiple Applications on Sensor Networks

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Junsung Kim, Karthik Lakshmanan, Raj Rajkumar

Macro-programming support

- **In-network programming**
- **Why?**
 - Usability
 - Lack of Technical Expertise with non CS people
 - Cost of re-program-ability
 - Faster deployment
 - Heterogeneous Hardware/Software
- **Consider a network with hundreds of nodes**
 - Few minutes per node can mean hours
- **More “Qualitative” advantages**

Sensor Network as Infrastructure



Occupancy Checking and Room Climate Control

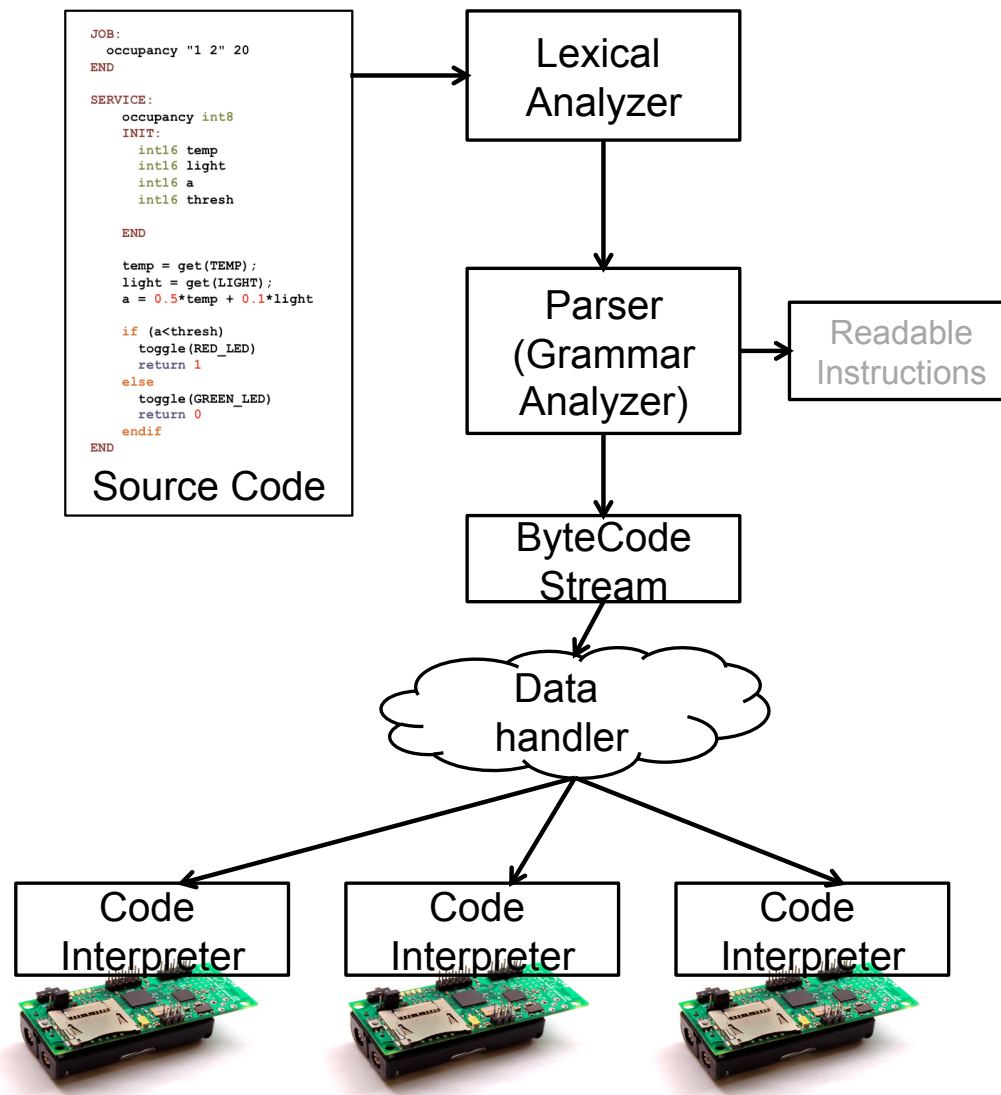
- Several applications on same sensor network
- Geographically distributed sensor network
- Limited flexibility and usability



Challenges for Concurrent applications

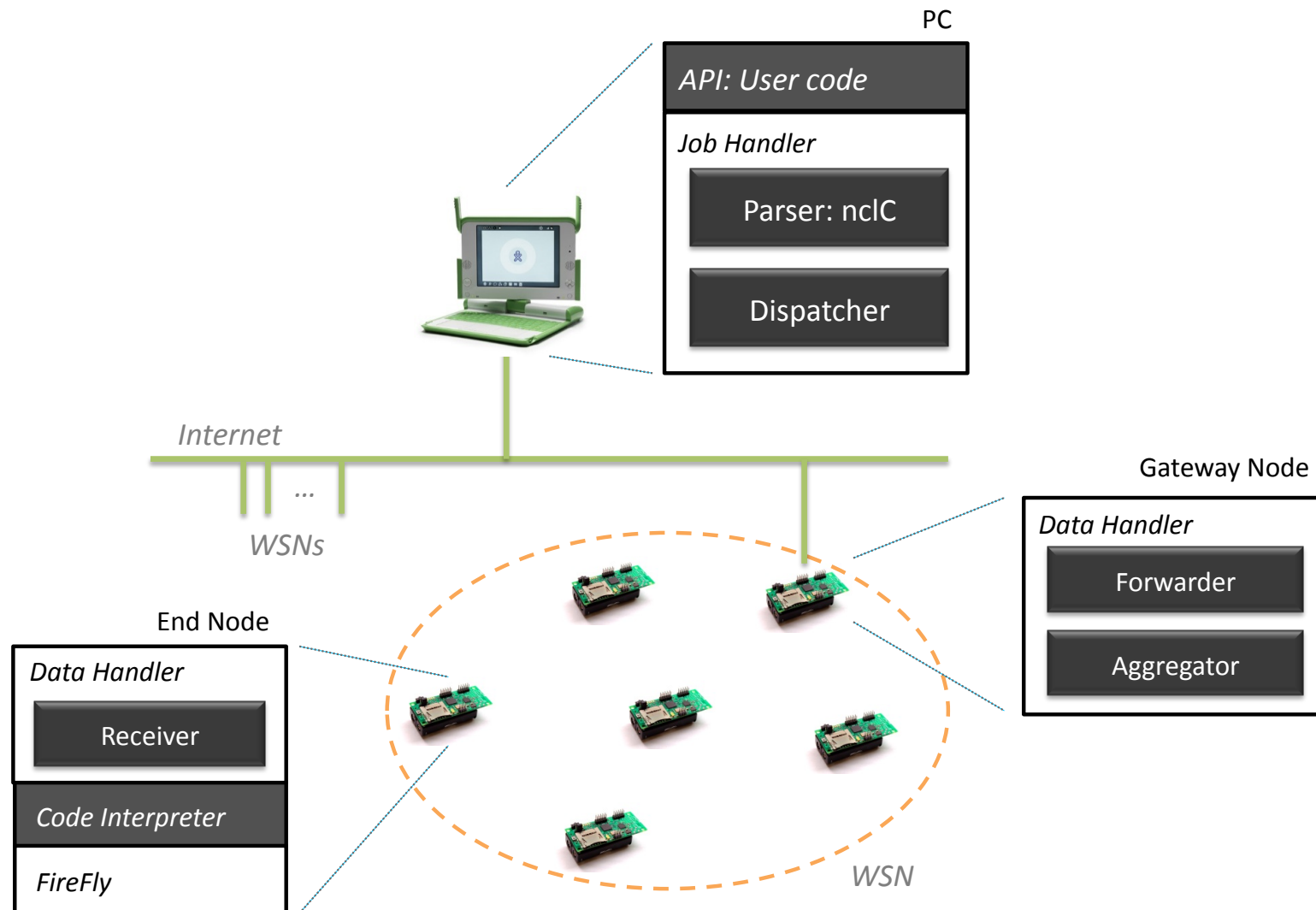
- **User Interface**
 - Database queries, virtual machine etc..
- **Operating System Support**
- **Packets through multiple applications over multi-hop network**
- **Data Aggregation**
- **Minimizing the overhead**
 - Frequency of Processor and Radio On/Off
 - Network flooding
 - Seamless backend handling
- **Tradeoff between Control and Abstraction**

Macro Programming System

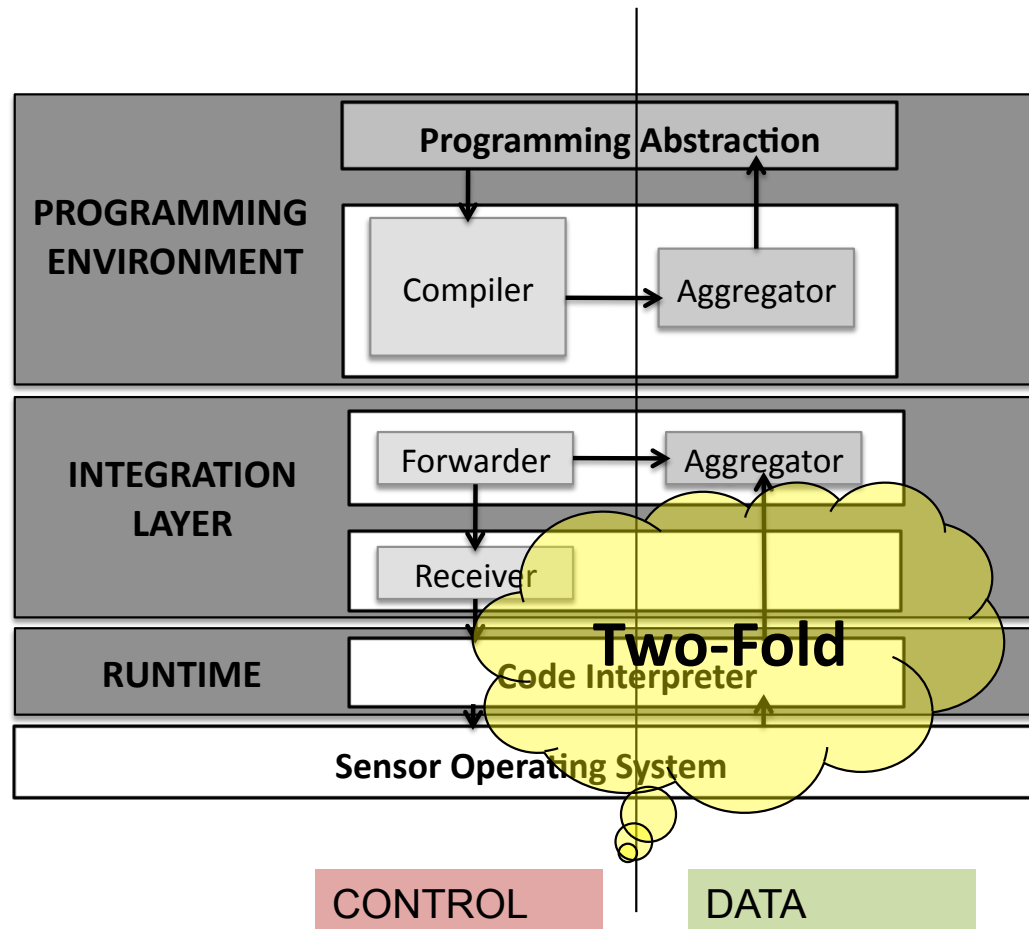


- nanoCL: Small Composable Language for Sensor Networks
- Abstracts away from lower-level details
- Supports various data-types and library functions
- Independent of the type of sensor nodes

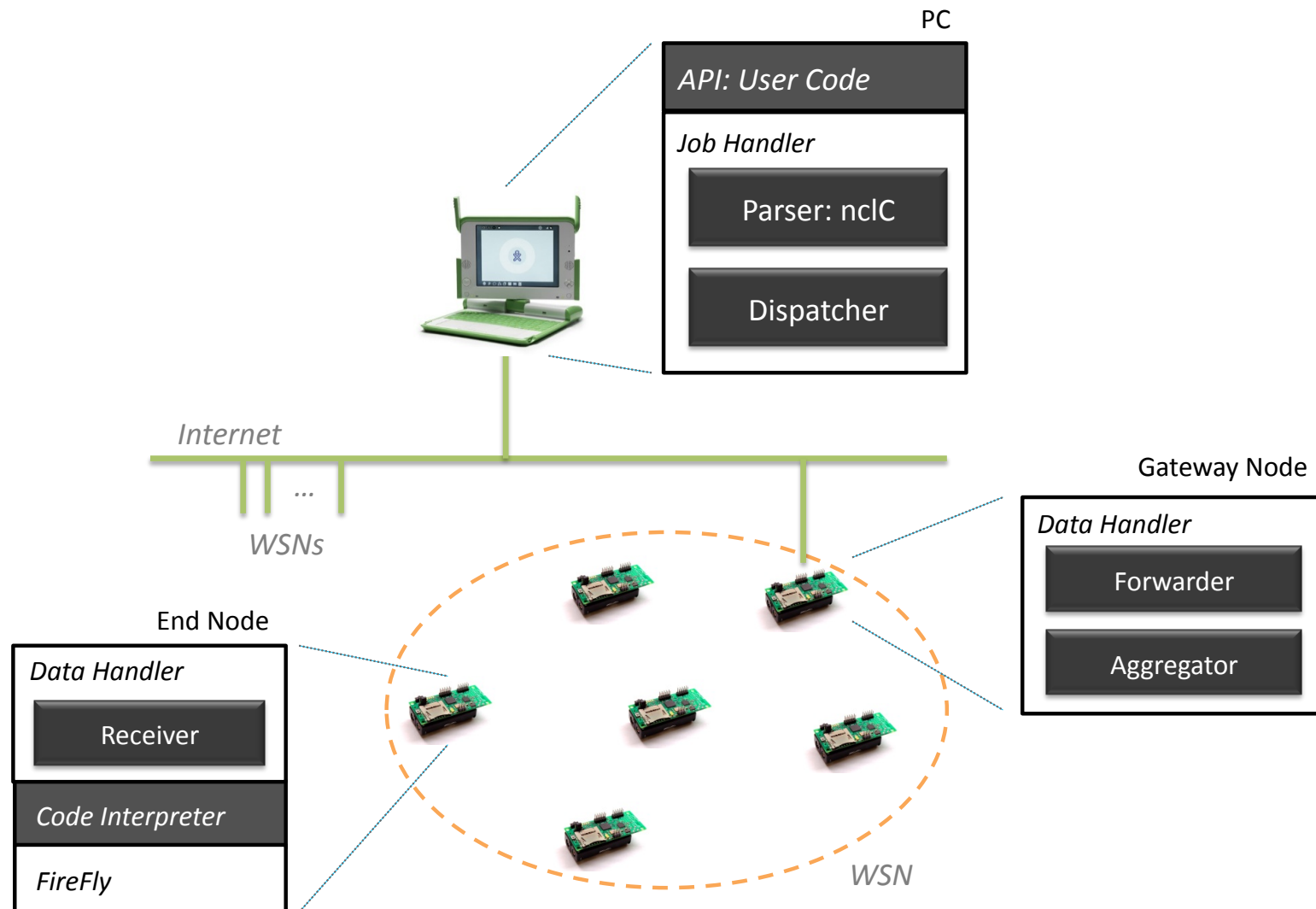
System Architecture



Control and Data Flow



System Architecture Outline

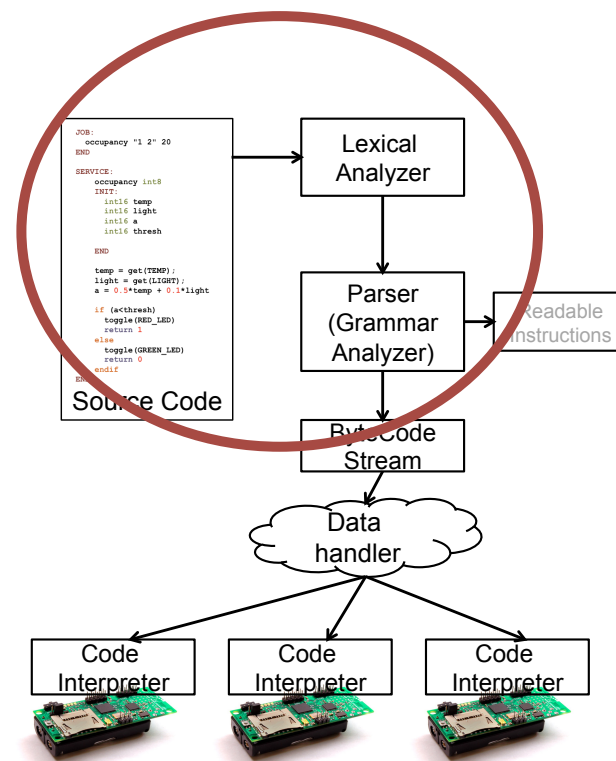


sMapReduce Programming Abstraction

- **Natural-fit to sensor network operation**
 - Map the “functionality” to sensor node
 - Gather the data through the network tree (Reduce)

- **Inspired from Google’s MapReduce**

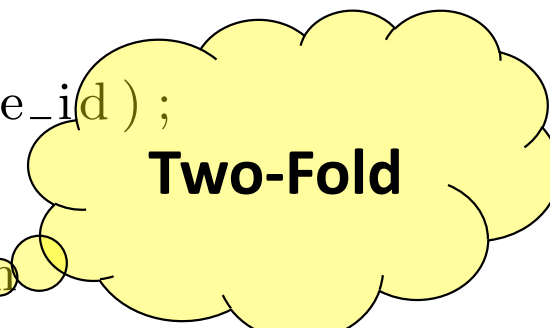
- **Key Features**
 - Balanced abstraction and control
 - Easy debugging
 - Two-fold operation



Simple Temperature Collection Example

```
1 smap(service_name , list_of_nodes , period) {  
2   for each node in list_of_nodes  
3     temp_value = gets(TEMP);  
4     smap_emit(temp_value , node_id);  
5   end
```

(a) sMap Function

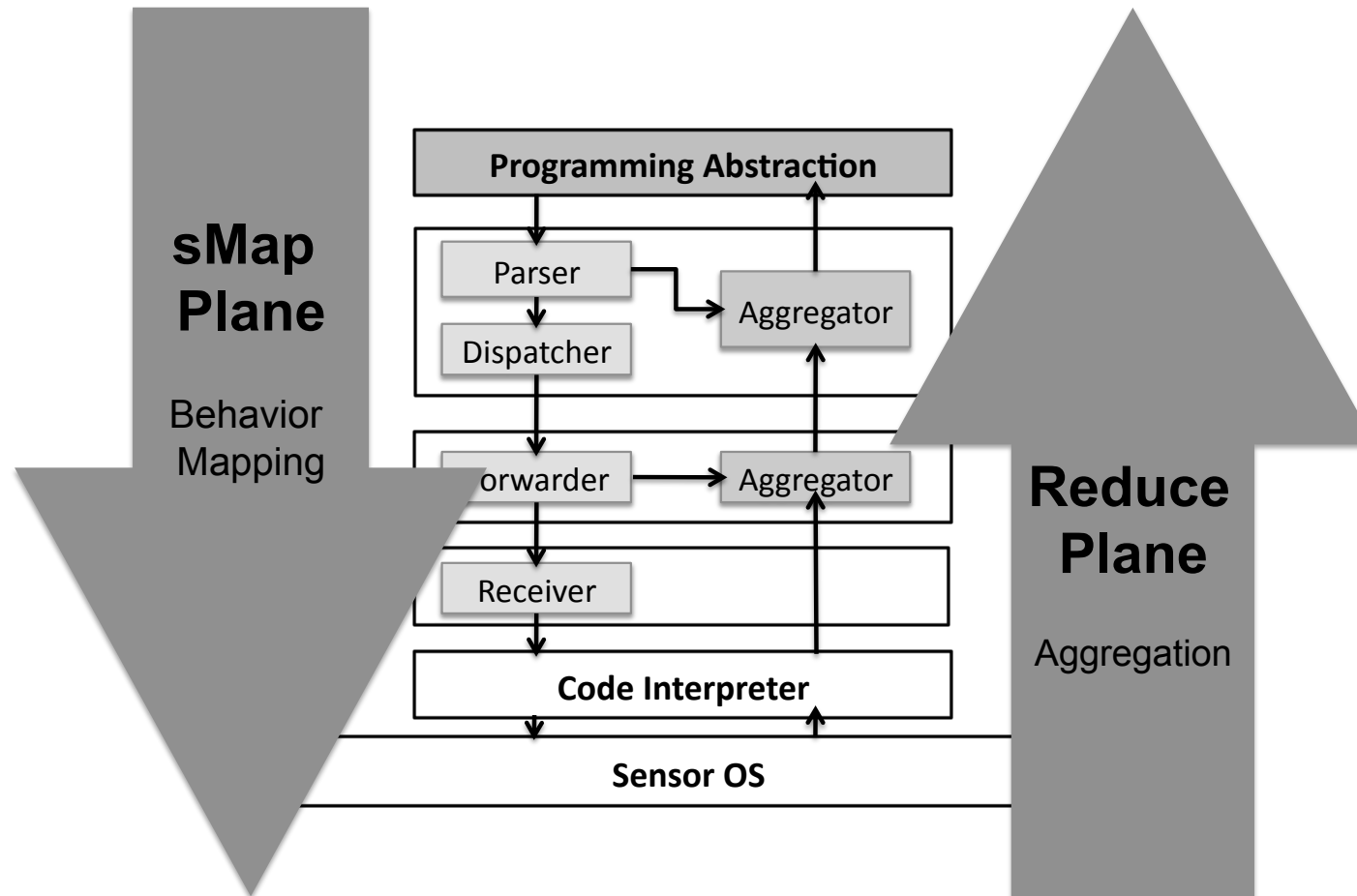


Two-Fold

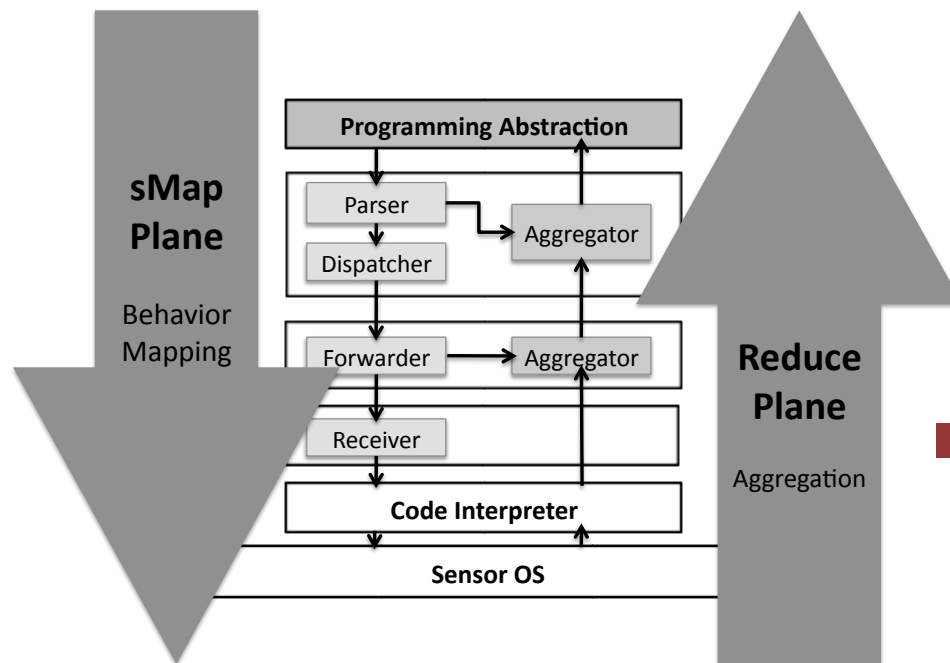
```
1 reduce(data , list_of_nodes) {  
2   for each node in INNER.list_of_nodes  
3     sum += data.temp_value; //AGGREGATION  
4   end  
5   return sum;  
6 }
```

(b) Reduce Function

Operation to UI Correlation



sMap and Reduce planes



- Left plane handles mapping the behavior
- Right plane handles the aggregation of data through the network

Target tracking application

```

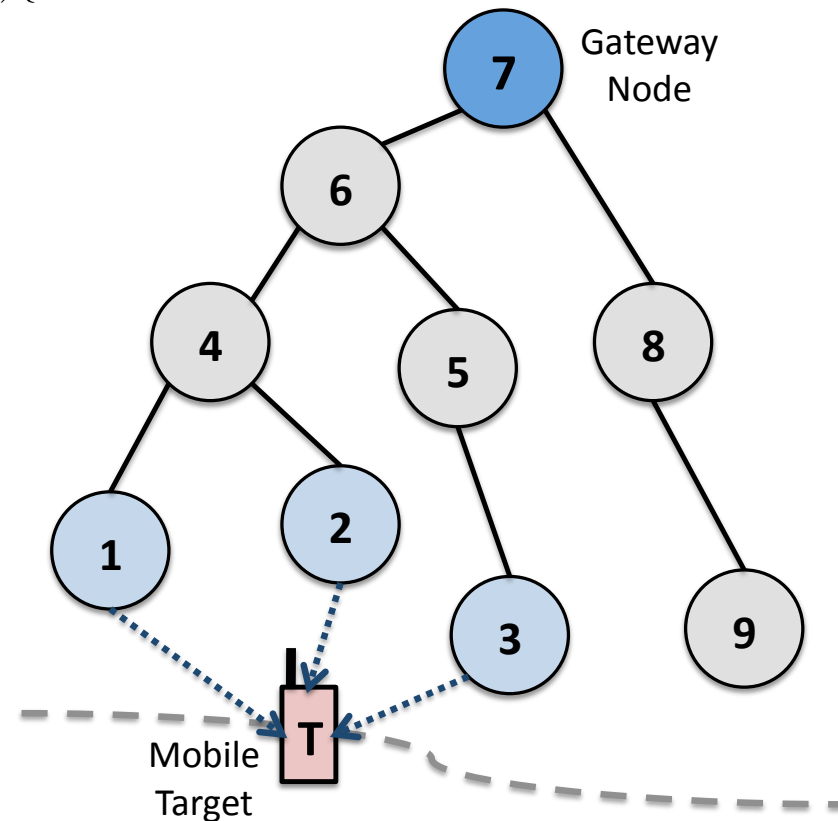
1 smap(target_track, list_of_nodes, period){
2   for each node in list_of_nodes
3     rssi_v = get(RSSI);
4     ts = get(time);
5     smap_emit(rssi_v, ts, node_id, loc);
6   end

```

```

1 reduce(data, list_of_nodes){
2   for each node in INNER.list_of_nodes
3     if(data.loc != NULL)
4       return data.loc;
5     else
6       if(max(ts)-min(ts) <= win
7         && size(data.rssi_v) >= 3)
8         triangulate(rssi_v, loc);
9       else
10        return data;
11      end
12    end
13  end
14 }

```



Write-ability

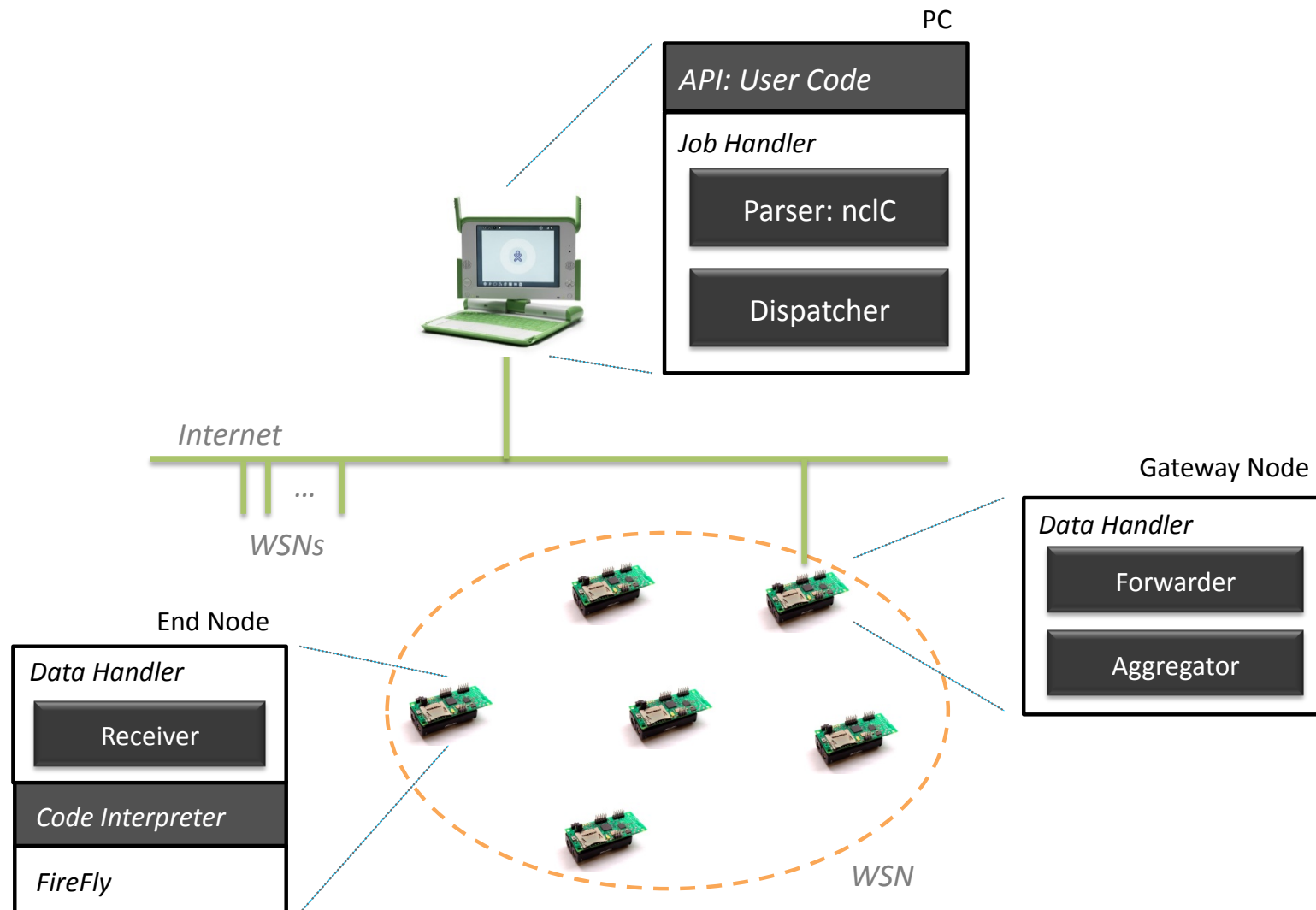
- **Simple C like syntax**
- **Library functions**
 - `gets ()` for accessing sensor data
 - Arithmetic operations
 - `int/uint` data types
 - `set ()`, `get ()`, `toggle ()` for GPIO pins
 - `for` and `while` loops
 - `if/else` constructs
 - `return` values to collect data

Source Lines of Code Comparison

Application	NanoCF	Operating System
Temperature Collection	5	80
Occupancy Monitoring	20	205
Target Tracking	20	~ 300 - 400

nanoCL Code	Corresponding Assembly	ByteCodes Generated by nanoCL Compiler
<pre> JOB: dummyservice "1 2 3 4 5" 100 MIN ENDJOB SERVICE: dummyservice int8 int8 INIT: int8 aa int8 bb int8 cc ENDINIT aa = gets(TEMP) bb = gets(LIGHT) clt (LED RED) cc = (bb/100) + (aa/100) if(cc > 15) set(LED RED) print(cc) endif wait(100) ENDSERVICE </pre>	<pre> No of Instructions: 35 SECTION INIT int8 a (aa) int8 b (bb) int8 c (cc) int16 d int16 e int16 f int16 g int16 h int16 i int16 j ENDINIT SECTION SERVICE GETS TEMP aa GETS LIGHT bb CLR LED RED AEQ d 100 DIV e b d AEQ f 100 DIV g a f ADD h e g MOV c h AEQ i 15 GT c i IF GOTO 11 LABEL 12 AEQ j 100 WAIT j ENDSERVICE REPEAT 0x00 0x64 LABEL 11 SET LED RED PRINT c GOTO 12 </pre>	<pre> 0x56, 0x58, 0xff, 0xff, 0x5c, 0x60, 0x61, 0xff, 0x5c, 0x60, 0x62, 0xff, 0x5c, 0x60, 0x63, 0xff, 0x5c, 0x61, 0x64, 0xff, 0x5c, 0x61, 0x65, 0xff, 0x5c, 0x61, 0x66, 0xff, 0x5c, 0x61, 0x67, 0xff, 0x5c, 0x61, 0x68, 0xff, 0x5c, 0x61, 0x69, 0xff, 0x5c, 0x61, 0x6a, 0xff, 0x59, 0xff, 0xff, 0xff, 0x56, 0x5a, 0xff, 0xff, 0x30, 0x61, 0x90, 0x00, 0x30, 0x62, 0x96, 0xff, 0x41, 0x95, 0x03, 0xff, 0x16, 0x64, 0x00, 0x64, 0x1a, 0x65, 0x62, 0x64, 0x16, 0x66, 0x00, 0x64, 0x1a, 0x67, 0x61, 0x66, 0x0d, 0x68, 0x65, 0x67, 0x17, 0x63, 0x68, 0xff, 0x16, 0x69, 0x00, 0x0f, 0x11, 0xff, 0x63, 0x69, 0x51, 0xff, 0xff, 0xff, 0x53, 0x11, 0xff, 0xff, 0x54, 0x12, 0xff, 0xff, 0x16, 0x6a, 0x00, 0x64, 0x44, 0x6a, 0xff, 0xff, 0x5b, 0xff, 0xff, 0xff, 0x45, 0xff, 0x00, 0x64, 0x54, 0x11, 0xff, 0xff, 0x40, 0x95, 0x03, 0xff, 0x31, 0x63, 0xff, 0xff, 0x53, 0x12, 0xff, 0xff, </pre>

System Architecture Outline



Data Handler Functions and Features

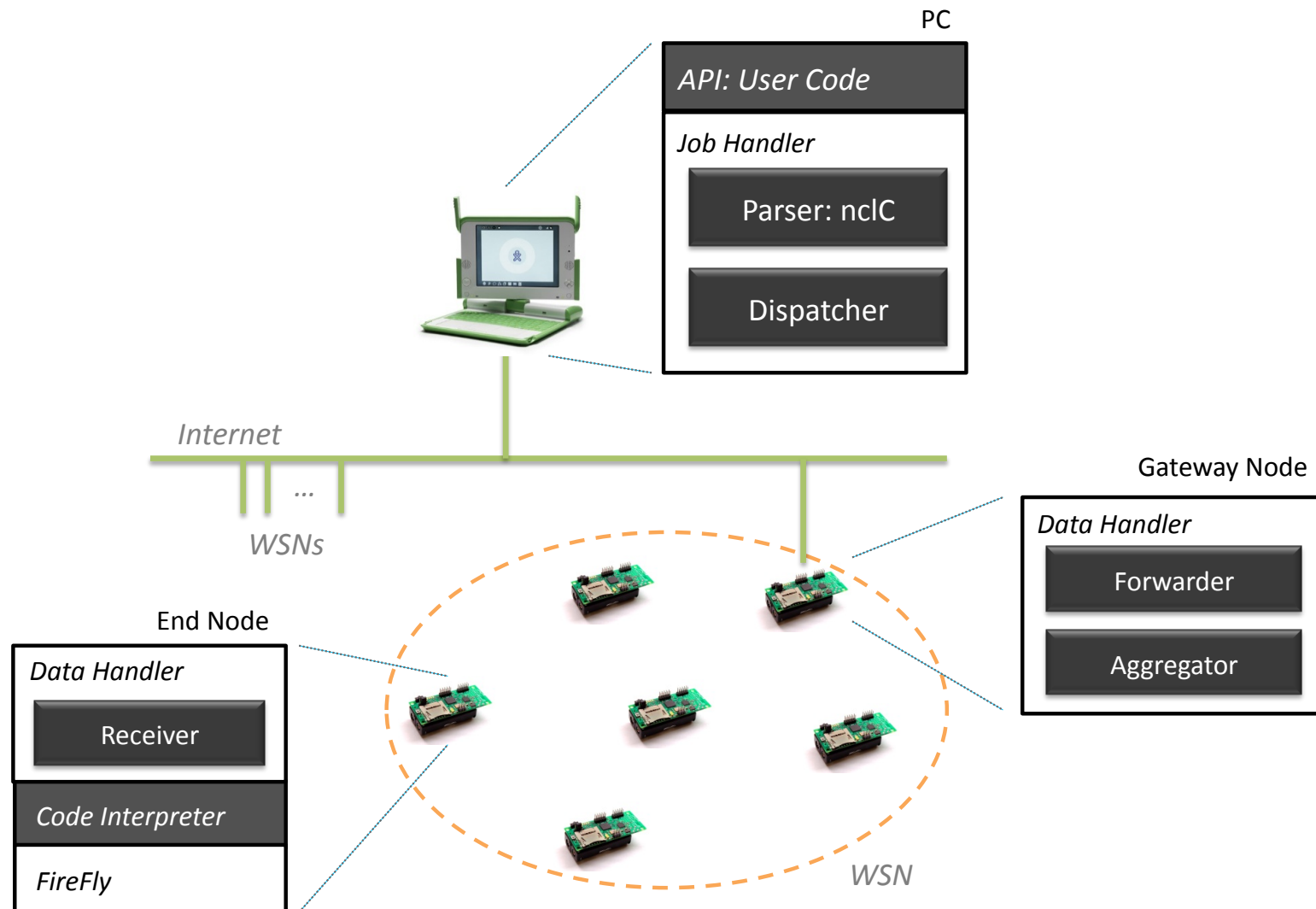
■ Functionalities

- Byte-code transfer
- Data transfer and aggregation
- Radio resource management

■ Features

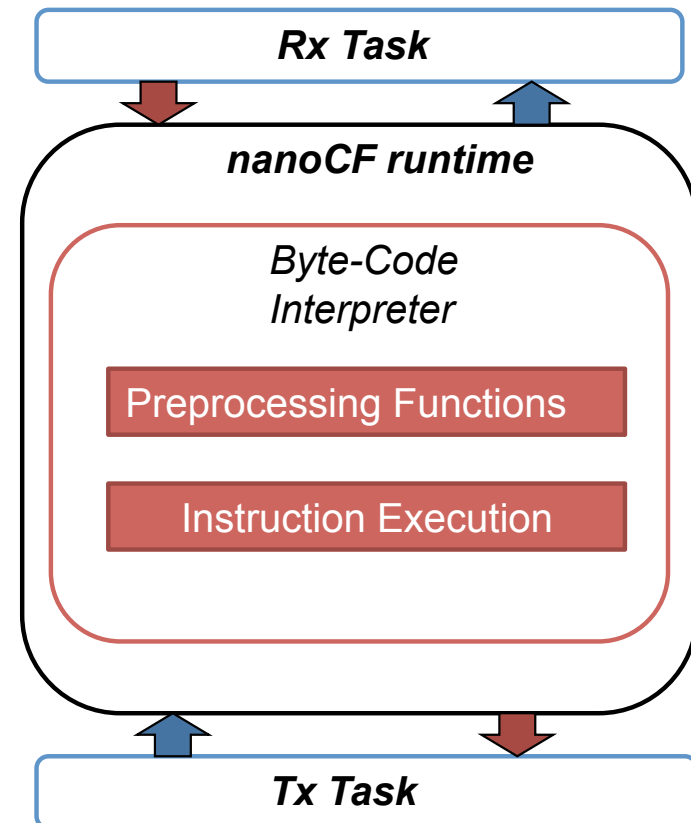
- Routing table management
- Fault-tolerant packet delivery
 - Retransmission
 - Random back-off delay between responses
- Application management
 - Tracking application transaction
 - Starting and terminating applications

System Architecture Outline



Code Interpreter

- Rx Task re-arranges received packets based on sequence
- Runtime pre-processes symbols and labels in the stack
- Interprets the instructions, evaluates values
- Sends the response value back to the gateway



Challenges for Concurrent applications

- User Interface
 - Database queries, virtual machine etc..
 - Operating System Support
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- **Minimizing the overhead**
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Task and Packet Scheduling

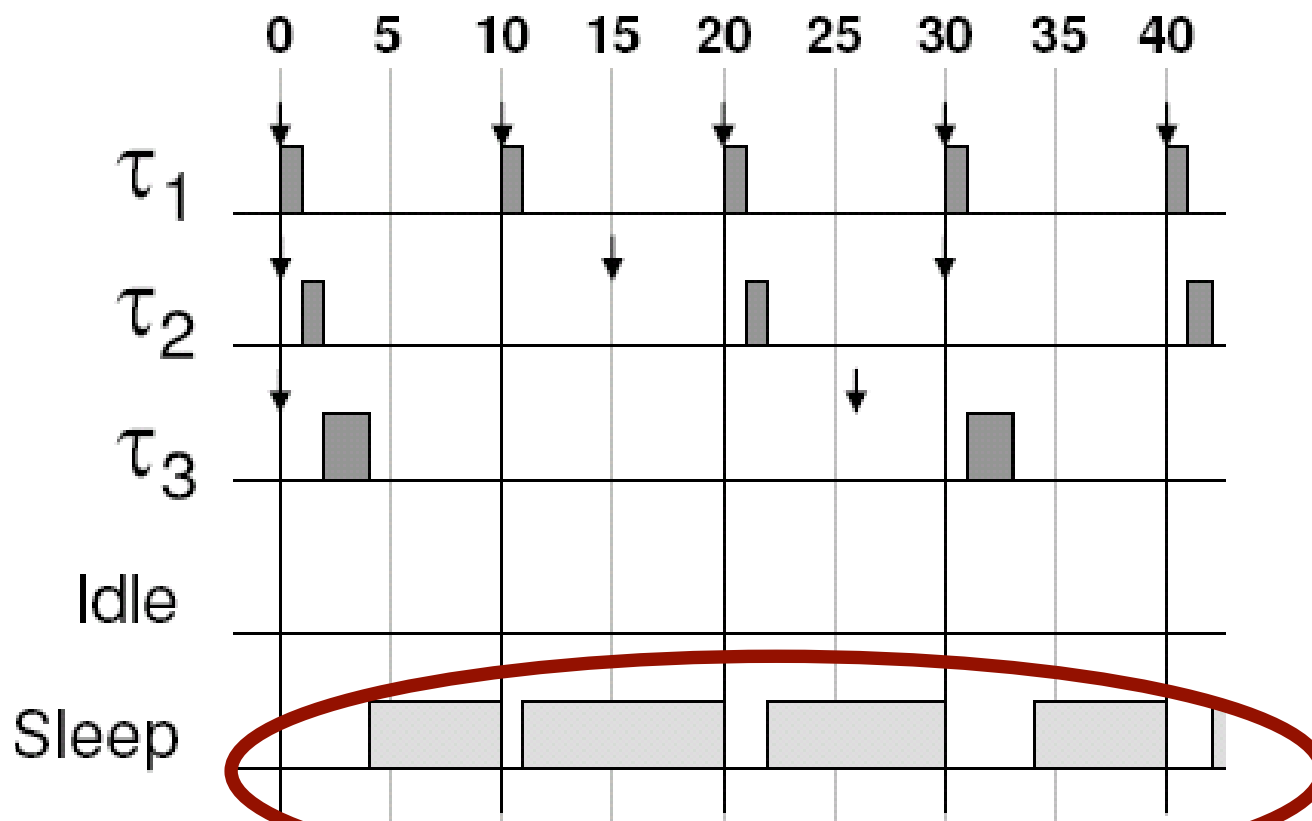
■ Typical Microprocessor operation states:

Power State	Power (mW)	Upward Transition Time
Active	30 mW	n/a
Idle	6 mW	6 μ s
Sleep	5 μ W	5 ms



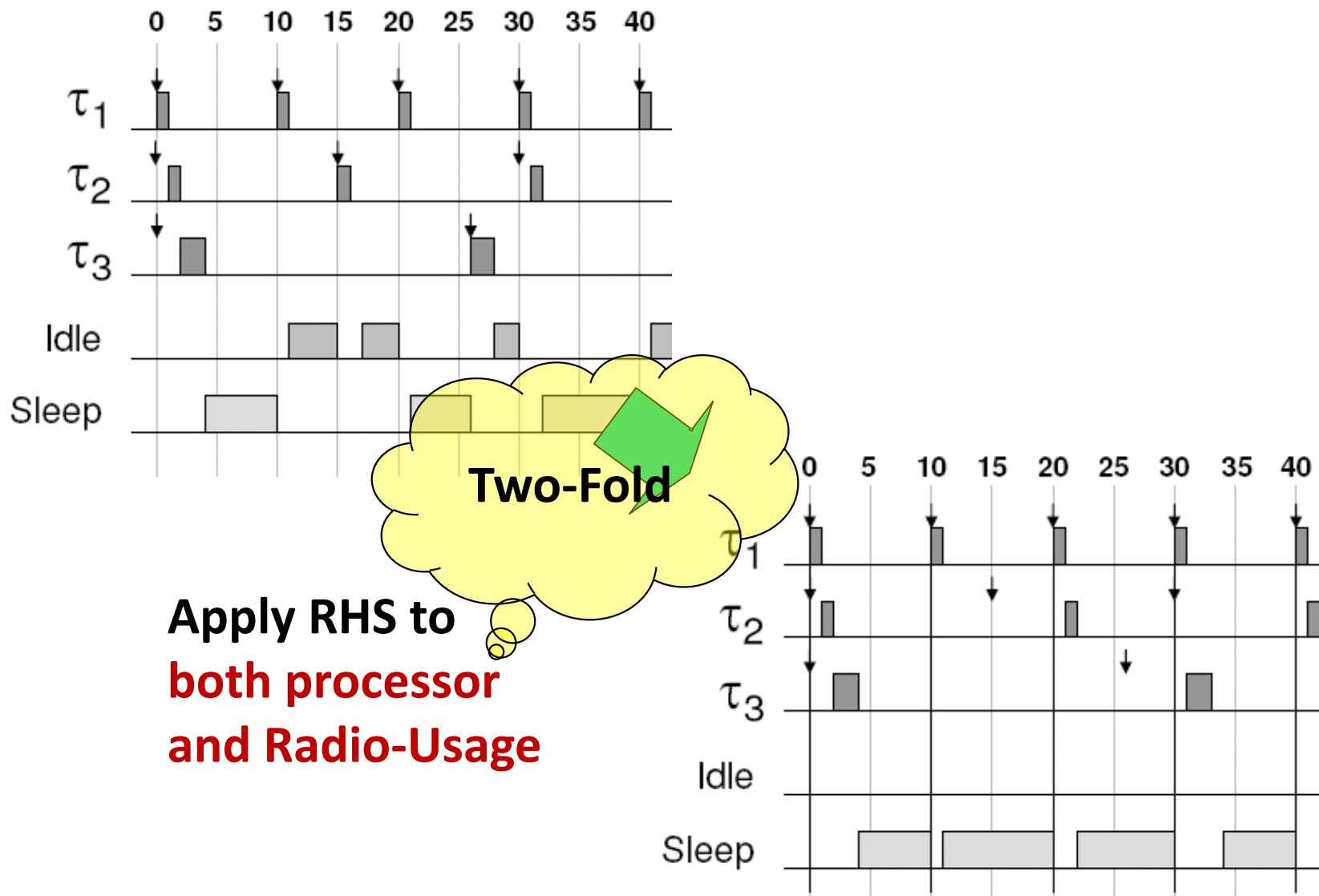
Rate Harmonized Scheduling¹

- Pick a harmonizing period (\leq shortest period)
- Release tasks only at the harmonizing interval



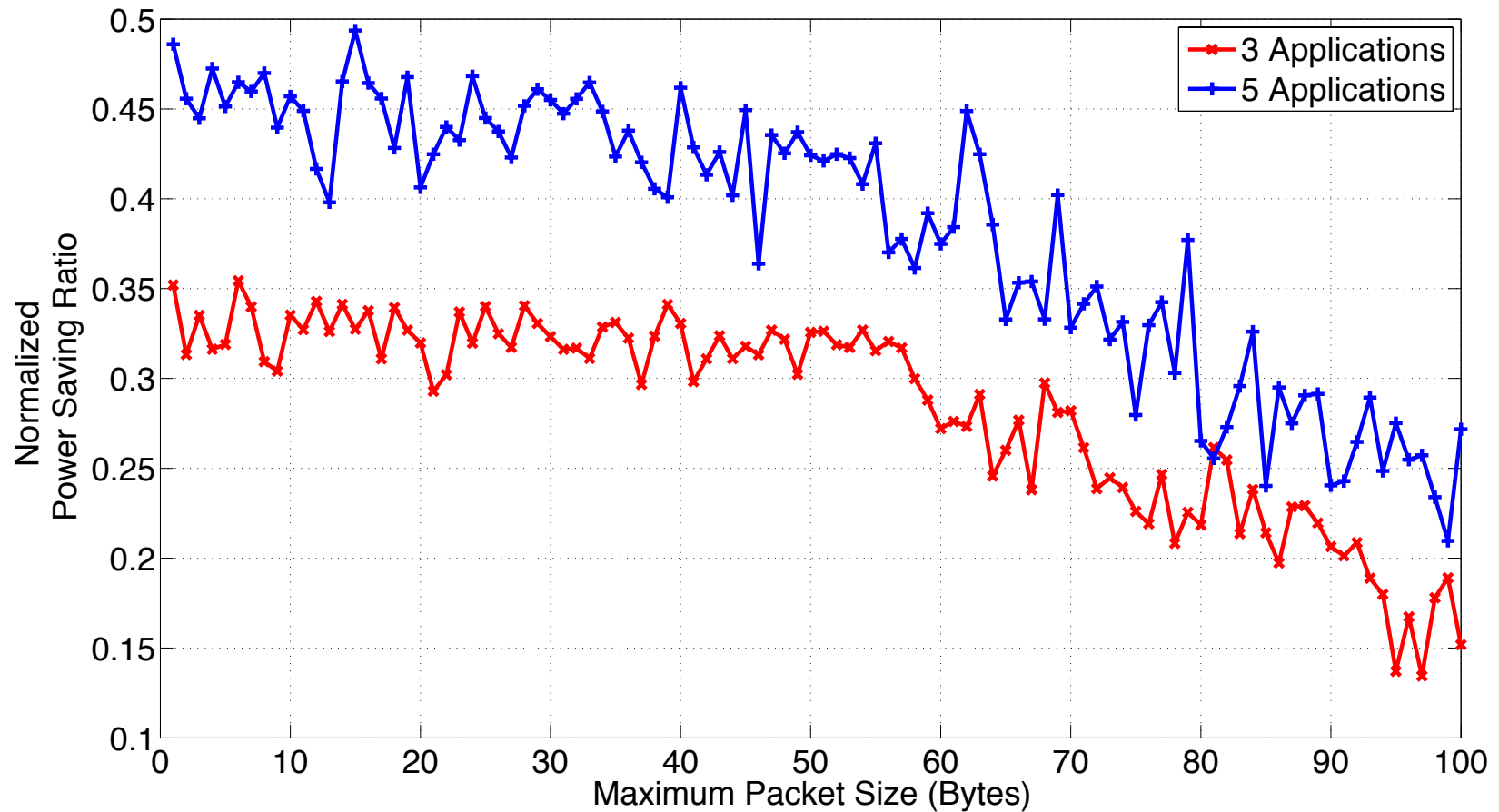
Anthony Rowe, Karthik Lakshmanan, Haifeng Zhu, Raj Rajkumar, ["Rate-Harmonized Scheduling for Saving Energy"](#). Proceedings of the 29th IEEE Real-Time Systems Symposium (RTSS), December 2008.

Transformation



Apply RHS to
both processor
and Radio-Usage

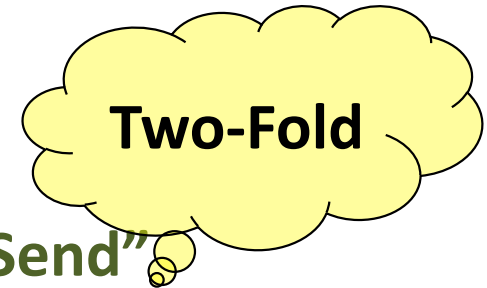
Power Saving in Radio



Not-So-Future Future-Work

- Optimize multiple applications
- Reduce the redundancy in applications
- Applications centered around “Sense & Send”
- Remove the double work of sensing
- Sending already addressed

- **Sample Light sensor only once**
 - Share data among multiple applications

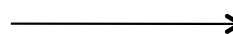


Longest Common Subsequence

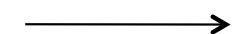
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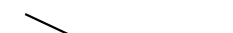
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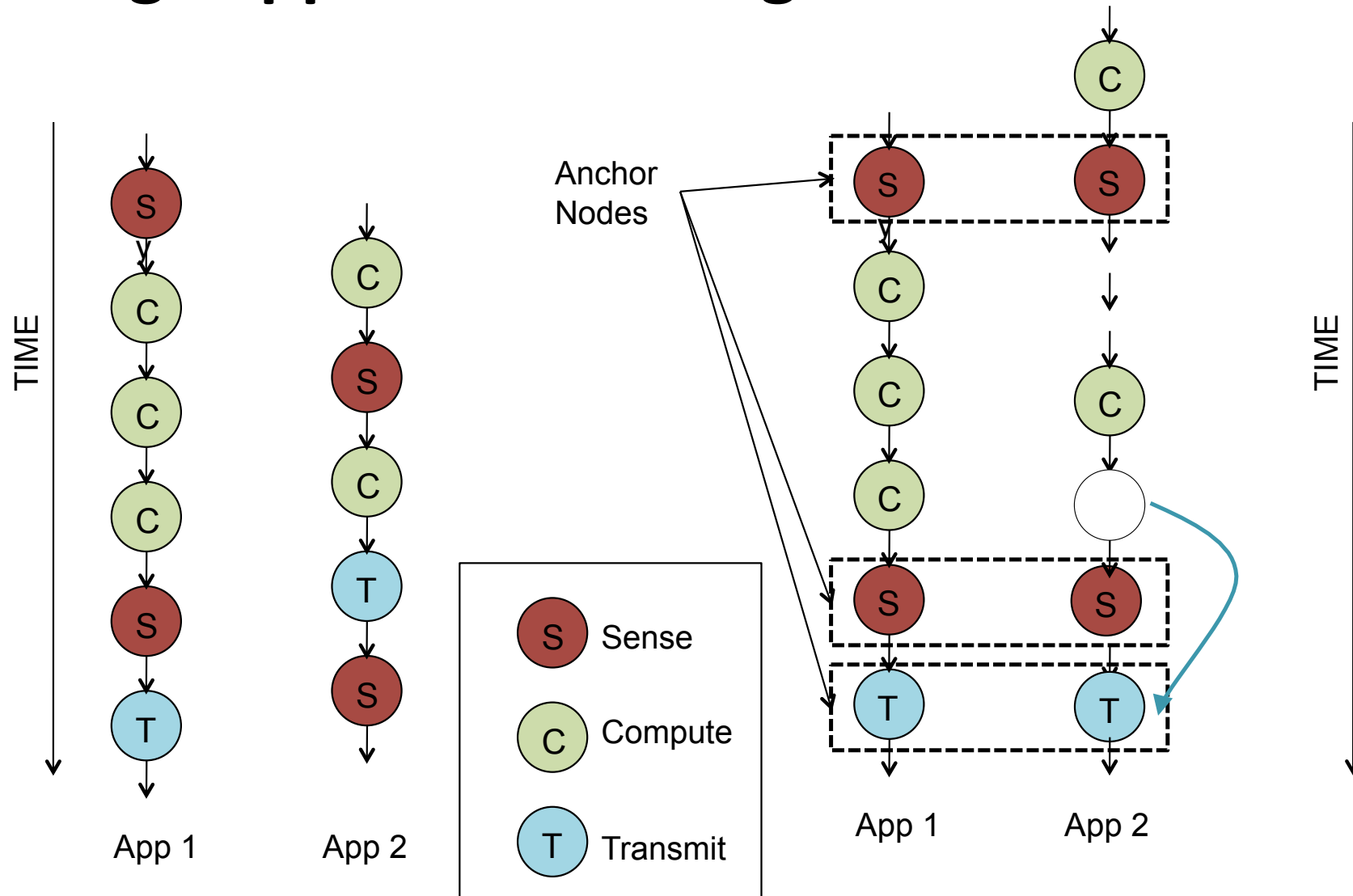
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Merge applications using LCS



Two-Fold

I could go on with more slides

BUT VIK STOPS HERE 😊